

# Complex Networks

## Problem Sheet 7

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**\*\* Please hand in solutions to questions 2 on this sheet \*\***

- Let  $K_4$  denote the complete graph on 4 vertices.
  - Draw  $K_4$ .
  - Compute exactly the expected number of copies of  $K_4$  in  $G(n, p)$ , the Erdős-Rényi random graph on  $n$  vertices where each edge is present with probability  $p$ , independent of the others.
  - Compute the variance of the number of copies of  $K_4$  in  $G(n, p)$  from first principles. It is enough if your answer gets the correct scaling in  $n$  and  $p$ . You can ignore any constants that don't depend on  $n$  or  $p$  in your calculations. You can also ignore terms in  $n$  and  $p$  that grow more slowly than the dominant term.
  - Find an  $\alpha_c \in (0, \infty)$  (or prove that none exists) such that the following is true:

$$\mathbb{P}(G(n, n^{-\alpha}) \text{ contains a copy of } K_4) \rightarrow \begin{cases} 0, & \text{if } \alpha > \alpha_c, \\ 1, & \text{if } \alpha < \alpha_c. \end{cases}$$

Justify your answer fully.

- Recall that  $G = (V, E)$  is bipartite if there exist vertex sets  $X$  and  $Y$  such that  $V = X \cup Y$ ,  $X \cap Y = \emptyset$  and  $E \subseteq X \times Y$ . In words,  $X$  and  $Y$  partition the vertex set, and there is no edge between two elements of  $X$  or two elements of  $Y$ .

The random bipartite graph  $G(n, n, p)$  has  $2n$  vertices which can be partitioned as  $V = X \cup Y$ ,  $X \cap Y = \emptyset$ , with  $|X| = |Y| = n$ . Moreover, each edge in  $X \times Y$  is present with probability  $p$ , independent of the others. There are no edges in  $X \times X$  or  $Y \times Y$ .

- Let  $K_{2,2}$  denote the complete bipartite graph on  $2+2$  vertices. Draw  $K_{2,2}$ .
- Compute exactly the expected number of copies of  $K_{2,2}$  in  $G(n, n, p)$ .
- Compute the variance of the number of copies of  $K_{2,2}$  in  $G(n, n, p)$  from first principles. It is enough if your answer gets the correct scaling in  $n$  and  $p$ . You can ignore any constants that don't depend on  $n$  or  $p$  in your calculations.
- Find an  $\alpha_c \in (0, \infty)$  (or prove that none exists) such that the following is true:

$$\mathbb{P}(G(n, n, n^{-\alpha}) \text{ contains a copy of } K_{2,2}) \rightarrow \begin{cases} 0, & \text{if } \alpha > \alpha_c, \\ 1, & \text{if } \alpha < \alpha_c. \end{cases}$$

Justify your answer fully.

3. Let  $S_k$  denote the star graph on  $k$  nodes, consisting of a hub and  $k - 1$  leaves.

(a) Show that  $S_k$  is a balanced graph.

(b) Using the results in your notes for balanced graphs, find a value  $\alpha_k \in (0, \infty)$  such that

$$\mathbb{P}(G(n, n^{-\alpha}) \text{ contains a copy of } S_k) \rightarrow \begin{cases} 0, & \text{if } \alpha > \alpha_k, \\ 1, & \text{if } \alpha < \alpha_k. \end{cases}$$

Clearly state the result you will use before using it.

(c) The chromatic number of a graph  $G$ , denoted  $\chi(G)$ , is defined as the minimum number of colours required to colour the nodes of the graph in such a way that no two nodes with an edge between them have the same colour. Show that  $\chi(G) \leq d_{\max} + 1$ , where  $d_{\max}$  denotes the maximum degree of all nodes in  $G$ .

*Hint.* Consider a greedy algorithm that goes through the nodes in arbitrary order, assigning each node a colour distinct from that of all its already coloured neighbours. Show that, if  $d_{\max} + 1$  colours are available, then the greedy algorithm will never get stuck.

(d) Provide either an upper or a lower bound on  $\chi(G)$  that holds with high probability for  $G$  drawn from  $G(n, p)$ . You should state your result precisely, and justify it fully.