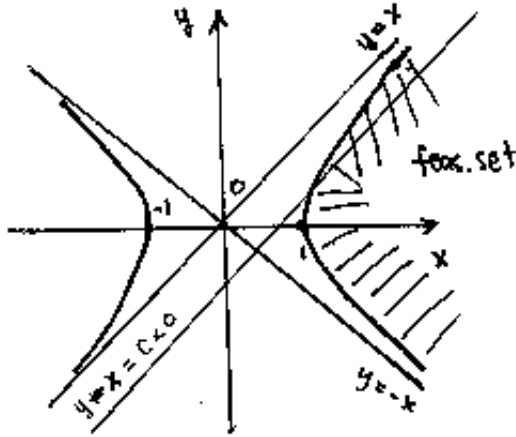


General optimisation problems

1. Clearly, $x^* = 0$, value 1.
2. One does not need a Simplex method to see that $(1, 1)$ is an optimal solution, value 2.
3. No optimal solution: as $(x, y) \rightarrow (1, 1)$, the value approaches 2, never reaching it.
4. Clearly when $x^2 + y^2 + z^2$ is minimal, i.e. $(0, 0, 0)$, value $e^0 = 1$. If Min is replaced by Max, then the problem has no optimal solution, because the feasible set is an ellipsoid (3D ellipse) with major axes $1/\sqrt{2}, 1/\sqrt{3}, 1/2$, the boundary not included, so the points $(\pm 1/\sqrt{2}, 0, 0)$ where $x^2 + y^2 + z^2$ achieves its supremum on the feasible set are not feasible.
5. The feasible set is a rhombus with vertices $(1, 0), (0, 1), (-1, 0), (0, -1)$, the objective is the distance from the origin. The maximum value 1 is attained at either vertex, as the rhombus is inscribed into the unit circle.
6. On the plane, consider a line $y - x = 0$ and a hyperbola $x^2 - y^2 = 1$, which asymptotically approaches it (see Fig.).



Then the feasible set is the right branch of the hyperbola and its interior, and so there is no optimal solution, as $y - x \rightarrow 0$ from below for a point on a hyperbola, going to infinity. Namely, a line $x - y = C$ for any $C < 0$ will end up entering the feasible set; however $y = x$ is still unfeasible.

Linear programming

1. Variables: x_1 - number of Grumpies, x_2 - number of Sleepies, x_3 - number of Bashfuls to be produced per week. Let $x = (x_1, x_2, x_3)$.

Then the LP is:

$$\max 4x_1 + 2x_2 + 5x_3, \text{ such that}$$

$$\begin{cases} x_1 + 2x_2 + x_3 \leq 40 & \text{Machine A constraint} \\ 3x_1 + \quad \quad + 2x_3 \leq 50 & \text{Machine B constraint} \\ x_1 + 4x_2 + \quad \leq 45 & \text{Machine C constraint} \end{cases}, x \geq 0.$$

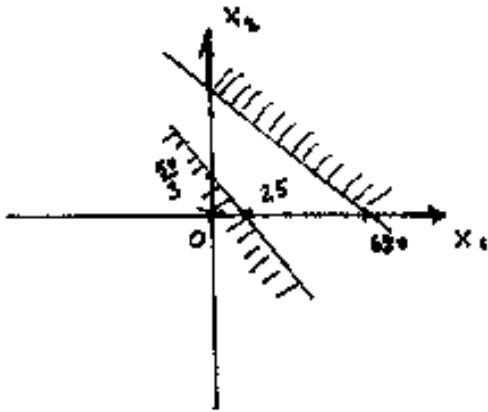
¹Please, let me know if you spot any errors.

Adding the extra equality constraint enables one to reduce the number of variables by one by expressing, e.g. $x_2 = 360 - x_1 - x_3$. Plugging into the objective and constraints yields:

max $720 + 2x_1 + 3x_3$, such that

$$\begin{cases} x_1 + x_3 \geq 680 \\ 3x_1 + 2x_3 \leq 50 \\ 3x_1 + 4x_3 \geq 1395 \end{cases}, x \geq 0.$$

The problem is clearly unfeasible, as one can see from the figure:

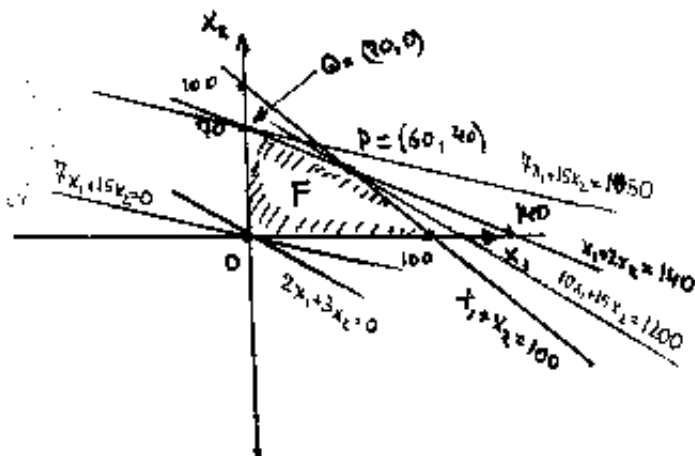


2. Variables: x_1 - number of Pasture acres, x_2 - number of Arable acres. Let $x = (x_1, x_2)$.

Then the LP is:

max $10x_1 + 15x_2$ (or $7x_1 + 15x_2$ in case of £7 per acre Pasture profit), such that

$$\begin{cases} x_1 + x_2 \leq 100 & \text{total area constraint} \\ 30x_1 + 60x_2 \leq 4200 & \text{total hours constraint} \end{cases}, x \geq 0.$$



By comparing the slopes of the lines on the figure, one can see that $P = (60, 40)$ corresponds to the optimal solution with value 1200 for the former profit function.

An extra acre changes the location of the optimal solution P to $(62, 39)$ and brings an extra value of 5.

For the second profit function, the optimal solution point "jumps" to $Q = (0, 70)$, with the value 1050. Note that if the profit function were $7x + 14x_2$, both P, Q and any point in between would yield an optimal value 980.

3. Unknowns in \mathbb{R}_+ : $e_1; e_2 \leq 1000, e_3 \leq 500$ and $r_1 \leq 100, r_2 \leq 130$ for the extracts to buy/remedies to mix.

The objective is now $\text{Max } 10r_1 + 13r_2 - 3e_1 - 4e_2 - 5e_3$, (profit from remedies' sales minus expenditure for extracts)

The constraints are such the amount of each of the three components extracted is sufficient to make the remedies: $.25r_1 + .2r_2 - .2e_1 - .3e_2 - .1e_3 \leq 0$,
 $.35r_1 + .1r_2 - .15e_1 - .3e_2 - .15e_3 \leq 0$,
 $.15r_1 + .3r_2 - .25e_1 - .45e_3 \leq 0$.

4. Introduce 16 unknowns $x_{ij} = \begin{cases} 1 & \text{if the corresponding candidate is chosen,} \\ 0 & \text{otherwise.} \end{cases}$

The objective is now to maximize $\sum_{i,j=1,\dots,4} e_{ij}x_{ij}$. The constraints are: $x_{ij} \geq 0$, for all i, j , as well as (i) for all $i = 1, \dots, 4, \sum_{j=1}^4 x_{ij} = 1$ - this takes care that no two team members do the same degree programme, as well as (ii) for all $j = 1, \dots, 4, \sum_{i=1}^4 x_{ij} = 1$ - this takes care that no university provides two team players.

There is a reasonable extra requirement that all x_{ij} be, in fact, integer, but it turns out to be superfluous - if the LP in question is solved via the simplex method to be learned soon, there will be no way to get non-integer values for the variables.

Linear algebra problems

1. While $\mathbf{u}^T \mathbf{u} = 1^2 + 2^2 + 3^2 = 14$ - the square of the Euclidean length of \mathbf{u} ,

$$\mathbf{u}\mathbf{u}^T = \begin{bmatrix} 1 & 2 & 3 \\ 2 & 4 & 6 \\ 3 & 6 & 9 \end{bmatrix}.$$

If A is $m \times n$, A^T is $n \times m$, so $A^T A$ is $n \times n$ and AA^T is $m \times m$, both symmetric. One can actually prove that they are both positive definite.

Furthermore,

$$AA^T = \begin{bmatrix} 14 & 32 \\ 32 & 77 \end{bmatrix}, A^T A = \begin{bmatrix} 17 & 22 & 27 \\ 22 & 29 & 36 \\ 27 & 36 & 45 \end{bmatrix}, BB^T = \begin{bmatrix} 2 & -5 \\ -5 & 25 \end{bmatrix},$$

$$B^T B = \begin{bmatrix} 1 & -1 \\ -1 & 26 \end{bmatrix}, BA = \begin{bmatrix} -3 & -3 & -3 \\ 20 & 25 & 30 \end{bmatrix}, B^2 = \begin{bmatrix} 1 & -6 \\ 0 & 25 \end{bmatrix};$$

the products AB and A^2 are not defined.

2. B is a 4×4 identity matrix, for the only change that $b_{31} = 1$, rather than 0. C is a 5×5 identity matrix, for the following changes: $c_{22} = c_{55} = 0$, instead of 1 and $c_{52} = c_{25} = 1$ instead of 0. That is

$$B = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}, C = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 & 0 \end{bmatrix}.$$

3. First system: $x = (1, 2, 3)$, unique solution;
 Second system: infinitely many solutions, e.g. $x = (4 - x_3 - 1.2x_4, 2 - x_3 - 1.4x_4, x_3, x_4)$, basic solution $x = (4, 2, 0, 0)$, or $x = (2 + x_2 + .2x_4, x_2, 2 - x_2 - 1.4x_4, x_4)$, basic solution $x = (2, 0, 2, 0)$;
 Third system: inconsistent, as adding the first two equations and subtracting the third one results in $0 = 1$.
4. When solving a single system of equations $Ax = b$, one does a succession of pivots in the extended matrix $[A|b]$, pivoting only the elements, which sit to the left of the vertical bar. No matter what b , the succession of pivots is the same, targeting eventually to get the identity to the left of the vertical bar:

$$[A | b] \rightarrow [\text{Id} | x].$$

As the pivots are being made, *each column is transformed independently*, which enables one to add as many columns to the right of the vertical bar as desired, doing the same eros, but with longer rows, the algorithm yielding:

$$[A | b^1 b^2 \dots] \rightarrow [\text{Id} | x^1 x^2 \dots],$$

where x^i solves a linear system $Ax^i = b^i$, $i = 1, 2, \dots$. As a matrix product, it can be compactly written as $AX = B$, the columns of X being x^1, x^2, \dots , while the columns of B are b^1, b^2, \dots . Indeed, in the example of the 3×3 matrix in question, it boils down to solving

$$\begin{bmatrix} 2 & 2 & 1 \\ 2 & -1 & 2 \\ 1 & -1 & 2 \end{bmatrix} \begin{bmatrix} x_1^1 & x_1^2 & x_1^3 \\ x_2^1 & x_2^2 & x_2^3 \\ x_3^1 & x_3^2 & x_3^3 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 1 \\ 0 & 0 & 1 \end{bmatrix},$$

or $AX = \text{Id}$, thus by definition $X = A^{-1}$. The Gauss-Jordan method yields:

$$\left[\begin{array}{ccc|ccc} 2 & 2 & 1 & 1 & 0 & 0 \\ 2 & -1 & 2 & 0 & 1 & 0 \\ 1 & -1 & 2 & 0 & 0 & 1 \end{array} \right] \rightarrow \left[\begin{array}{ccc|ccc} 1 & 0 & 0 & 0 & 1 & -1 \\ 0 & 1 & 0 & .4 & -.6 & .4 \\ 0 & 0 & 1 & .2 & -.8 & 1.2 \end{array} \right].$$

Finally, $\det A = -5$, as a calculation shows. Note that for a 3×3 matrix, one can compute the determinant by the template on the Fig. 1.

$$\det A = \begin{bmatrix} * & * & * \\ * & * & * \\ * & * & * \end{bmatrix} - \begin{bmatrix} * & * & * \\ * & * & * \\ * & * & * \end{bmatrix} = a_{11} a_{22} a_{33} + a_{12} a_{23} a_{31} + a_{13} a_{21} a_{32} - a_{11} a_{23} a_{32} - a_{12} a_{21} a_{33} - a_{13} a_{22} a_{31}$$

Figure 1: How to compute the determinant of a 3×3 matrix

One adds products of elements in triples, connected by fat lines and subtracts products of elements in triples, connected by thin lines, e.g. for the matrix A in question

$$\det A = 2 * (-1) * 2 + 2 * 2 * 1 + 2 * (-1) * 1 - 1 * (-1) * 1 - 2 * 2 * (-1) - 2 * 2 * 2 = -5.$$

5. $\text{rank } A = 2$, because $\det A = 0$, however it contains a non-degenerate 2×2 submatrix, for instance

$$A = \begin{bmatrix} 1 & 2 \\ 4 & 5 \end{bmatrix},$$

whose determinant is nonzero.

6. Let $C = AB$. Let the elements of A be a_{ij} , the elements of A^T be $a_{ij}^* = a_{ji}$, the elements of B be b_{ij} , the elements of b^T be $b_{ij}^* = b_{ji}$, the elements of C be c_{ij} , the elements of $C^T = (AB)^T$ be $c_{ij}^* = c_{ji}$. By the multiplication rule

$$c_{ij}^* = c_{ji} = \sum_k a_{jk} b_{ki} = \sum_k b_{ki} a_{jk} = \sum_k b_{ik}^* a_{kj}^*,$$

but the right-hand side is nothing but the multiplication rule for the product $B^T A^T$.

For the inverses:

$$(B^{-1}A^{-1})AB = B^{-1}(A^{-1}A)B = B^{-1}B = \text{Id},$$

so by definition of an inverse, the matrix $(B^{-1}A^{-1})$ is an inverse for the matrix AB , that is $(AB)^{-1}$.

7. Rows \rightarrow Columns: if the rows of an $n \times n$ matrix A are linearly dependent, there exists a sequence of eros, which being applied to A produce a matrix B , whose last row is zero. Each single ero consists in multiplying A from the left by a very simple and non-degenerate matrix (Problem 2). Thus $CA = B$ for some non-degenerate matrix C (i.e. such that its inverse C^{-1} exists). In B the zero row can be omitted, whereupon there remains some matrix \hat{A} with n columns and $n - 1$ rows. So, the columns of \hat{A} are vectors in \mathbb{R}^{n-1} , and their number is n . But more than $n - 1$ vectors in \mathbb{R}^{n-1} are always linearly dependent, by definition of dimension!

$$\begin{bmatrix} A \\ \vdots \\ 0 \end{bmatrix} \rightarrow \begin{bmatrix} \hat{A} \\ 0 \dots 0 \end{bmatrix} = B$$

Thus the columns of B are also linearly dependent. But $A = C^{-1}B$, hence $\mathbf{a}^i = C^{-1}\mathbf{b}^i$, for the columns of A and B respectively. Then, because the columns of B are linearly dependent, so are the columns of A . Namely, if $\lambda_1\mathbf{b}^1 + \dots + \lambda_n\mathbf{b}^n = 0$ for some array of numbers $(\lambda_1, \dots, \lambda_n) \neq (0, \dots, 0)$, then $\lambda_1\mathbf{a}^1 + \dots + \lambda_n\mathbf{a}^n = C(\lambda_1\mathbf{b}^1 + \dots + \lambda_n\mathbf{b}^n) = 0$.

An argument Columns \rightarrow Rows follows by transposition.