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QUANTUM COMPUTATION Exercise sheet 3

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1. The QFT and periodicity.

- (a) Multiply out the matrices corresponding to the gates in the circuit for the quantum Fourier transform Q_4 , in the computational basis, and check that the result is what you expect.
- (b) Write the state $Q_8|3\rangle$ as a tensor product of three single-qubit states, each of the form $\frac{1}{\sqrt{2}}(|0\rangle + e^{2\pi i z}|1\rangle)$ for some binary fraction z (i.e. something of the form $(.x_{j-1}...x_0)$). Expand out the resulting state and check that the answer is what you expect.
- (c) Let $f : \mathbb{Z}_{16} \to \mathbb{Z}_4$ be the periodic function such that f(0) = 2, f(1) = 1, f(2) = 3, f(3) = 0, and f(x) = f(x-4) for all x (so f(4) = 2, etc.).
 - i. Work through all the steps of the periodicity determination algorithm, writing down the state at each stage, and assuming that the measurement outcome in step 3 is 1, and the measurement outcome in step 5 is 12. Does the algorithm succeed?
 - ii. Now assume that the measurement outcome in step 5 is 8. Does the algorithm succeed?

2. Shor's algorithm.

- (a) Suppose we would like to factorise N = 85 and we choose a = 3, which is coprime to N. Follow steps 3-5 of the integer factorisation algorithm to factorise 85 using this value of a (calculating the order of a classically!). You might like to use a computer.
- (b) Imagine we want to factorise N = 21 and we choose a = 4. Does the integer factorisation algorithm work or not?
- 3. Approximate implementation of the QFT (optional). This part proves a claim made at the end of Section 4 of the lecture notes. Define the distance D(U, V) between unitary operators U and V as the maximum over all states $|\psi\rangle$ of $||U|\psi\rangle V|\psi\rangle||$.
 - (a) Show that $D(\cdot, \cdot)$ is subadditive: $D(U_1U_2, V_1V_2) \le D(U_1, V_1) + D(U_2, V_2)$.

- (b) Show that $D(R_d, I) = O(2^{-d})$ and argue that the same holds for controlled- R_d .
- (c) Describe how to produce a quantum circuit for an operator \widetilde{Q}_{2^n} on n qubits such that \widetilde{Q}_{2^n} uses $O(n \log n)$ gates and $D(\widetilde{Q}_{2^n}, Q_{2^n}) = O(1/n)$.